

Expediting the transport of Data Center Flows (DAQ: Deadline-Aware Queue)

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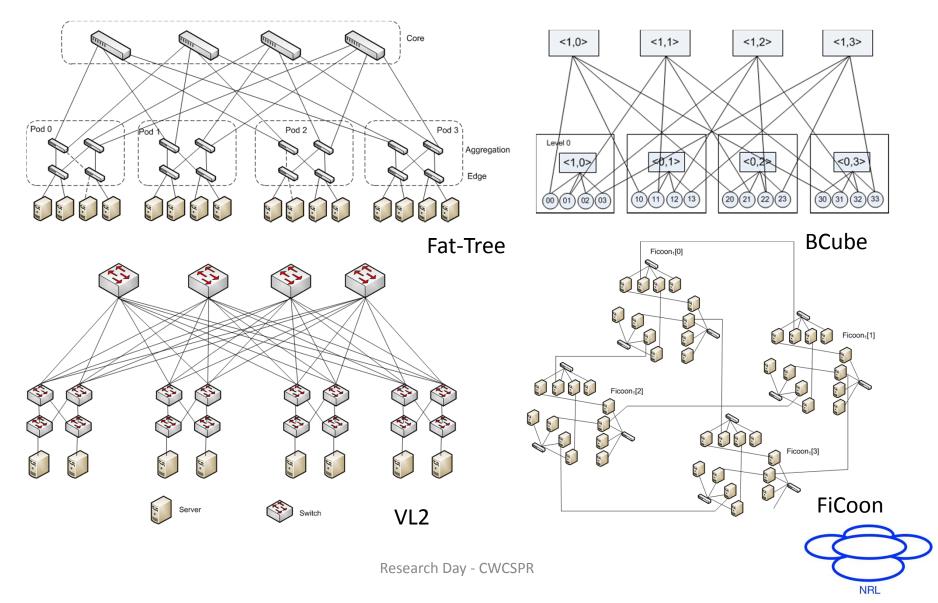
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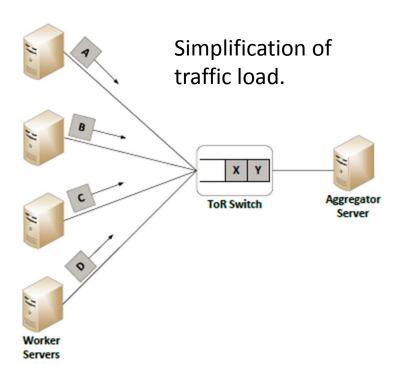




Examples of DC topologies



What is unique in Data Center Traffic? Partition-Aggregate Model







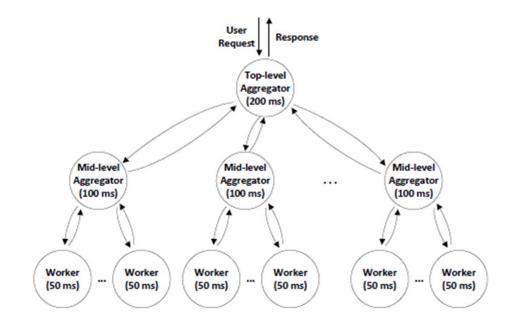
Data Aggregation

Flows may be associated with response deadlines

Deadlines are inherited by partial processes

For all flows, short Flow Completion times (FCTs) are desirable

For deadline-sensitive flows, short **Application Throughput** is desirable.



Data aggregation → Connection-Oriented Transport → Transmission Control Protocol (TCP)



Expected requirements of a Data Center (DC) Transport Protocol

- ➤ Maximize the number of flows completing transmission before deadlines
- Guarantee a high throughput for long flows.
- > Allow high, if not 100%, link utilization.
- > Achieve lossless transmissions.
- ➤ Minimize the amount of state information at switches

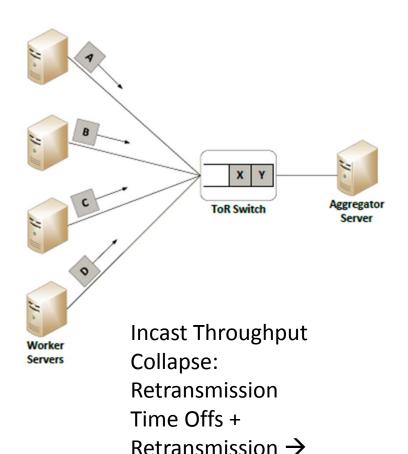




Why TCP is not good enough?

- Data Center Flows: Long + Short Flows
- Congestion
- Multiple flows concur at aggregation switches
- Lack of a centralized scheduler

Flow control mechanisms are not transmission speed aware → Long FCTs!



choke bandwidth





Existing Solutions

- Earlier Congestion Notification (ECN): DCTCP
- Rate Control: D2TCP, D3, PDQ (deadline aware)
- Congestion Control: RCP
- Pacing Schemes: HULL
- Load Balancing Schemes: DeTail, CONGA, RepFlow
- Switch Modification: DAQ





Deadline-Applicable Schemes

- RCP [Dukkipati05] assigns rate according to available bandwidth. Parameters must be tuned.
- DCTCP [Alizadeh10]: ECN + congestion window modification. Agnostic to deadlines.
- **D**³ [Wilson11] reserves transmission rates FCFS.
- PDQ [Hong12]: selects flows → earliest deadline first (EDF) and the shortest job first (SJF). High complexity.



Proposed Scheme: Deadline Aware Queue (DAQ) at DC Switches

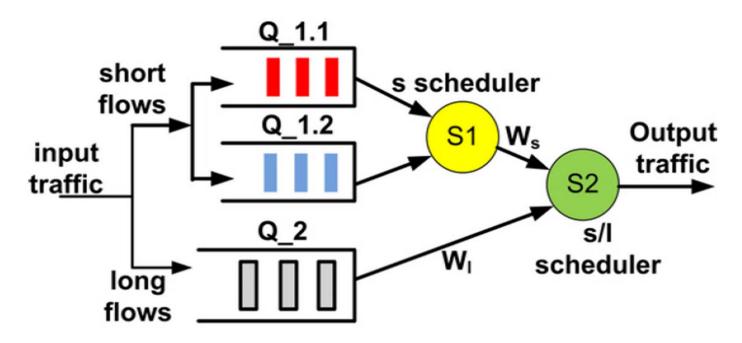
Objectives:

- Maximize application throughput
- Ensure minimum bandwidth for long flows
- Minimize flow-state information at switches
- Minimize modification to layered protocols





Switch Architecture

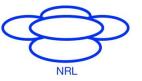


Use Three Queues: Urgent, Non-urgent, Long

Short flows: Urgent or Non-urgent

Long flows: long-flow queue + service

weighted scheduling

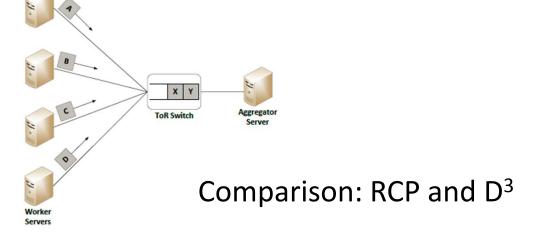


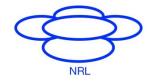


Test setup

- Loss-less flow control between
 - Senders and switch
 - Switch and receiver (aggregator)
- Large congestion window size instead of slow

start

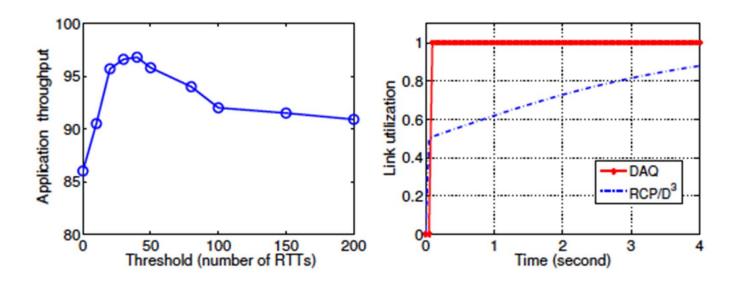






Impact of Urgent Threshold Value

Application throughput: No. on-time flows/All arrived flows

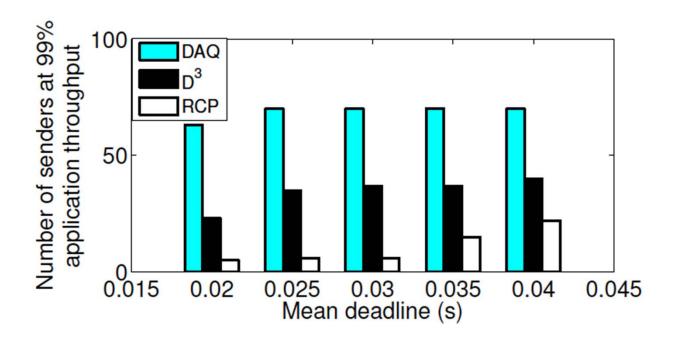


Flow size: 30KB, rate: 3600 flows/s Number of long flows: 5

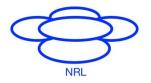




Supported number of senders

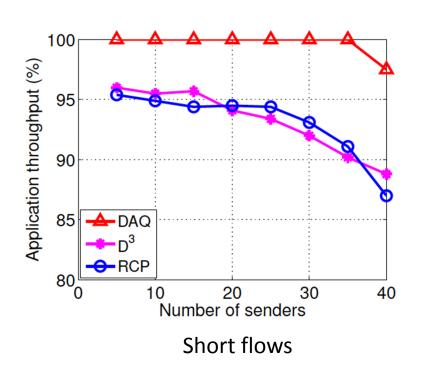


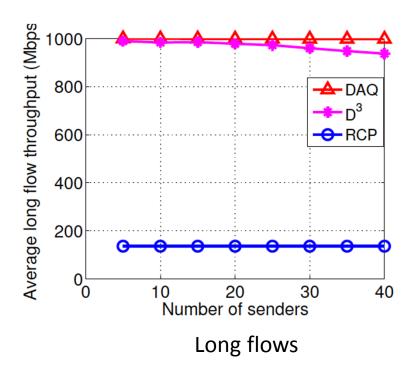
Number of concurrent senders for achieving 99% application throughput with flow size mean of 10 Kbytes and deadlines [20, 40] ms.





Application and Average Throughput





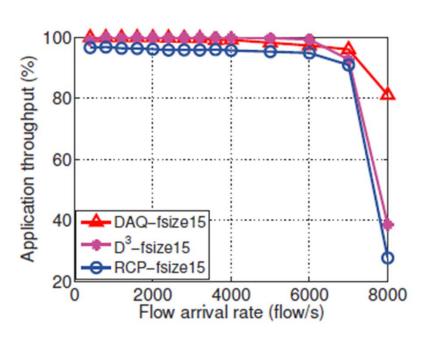
Short flow size: 15 Kbyte, long flow size: 100Mbyte (2).

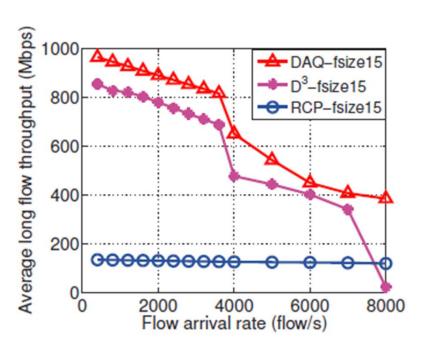
Short flow load: 0.3 % No. of senders: [5, 40]

Research Day - CWCSPR



Performance under short and long flows





(a) Application throughput for (b) Average long flows throughshort flows. put.

Short flow size: 15KB





Conclusions

- Deadline-oriented approach with small modification to transport layer.
- Urgent flows receive preferential service.
- Few urgent flows speedup transmission.
- DAQ achieves high Application Throughput
- Long flows receive minimum throughput through Weighted Round-Robin





Thank you

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