

# Coding for Leveraging Network Gains in 5G

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#### The Zettabyte Area



[Cisco Systems Inc.: The zettabyte area. White paper, 2015]

#### How Can Cellular Systems Keep Up?



[Nokia Networks: Looking ahead to 5G. White paper, April 2014]

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  - cooperative and opportunistic communication
  - improved multiple access techniques

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In this talk: How to leverage network gains for both error correction and compression with modern graph based codes?

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  - Multi-terminal coding and decoding (i.e., for relaying, cooperation, broadcast)

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- Partitioning into subcodes  $C_{\ell}$ ,  $\ell = 1, 2, ..., M$
- Can be seen as structured linear binning schemes
- Finite-field version of physical layer superposition codes

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- Interference channels
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#### Interference channels

- e.g., scheme from [Han & Kobayashi 1981], 2-user channel: message split up in public and private part (codes C<sub>1</sub> and C<sub>2</sub>)
- Many more: relay channels, cooperative diversity, wiretap channels, …

- Consider k by n generator matrix **G** of linear code C
- *M* information words  $\mathbf{i}_k$
- **Type-1** codes: Partitioning of **G** into subcodes  $C_{\ell}$  with generator **G**<sub> $\ell$ </sub> and rate  $R_{\ell} = k_{\ell}/n$

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$$\mathbf{c} = [\mathbf{i}_1, \, \mathbf{i}_2, \, \dots, \, \mathbf{i}_M] \, \mathbf{G} = [\mathbf{i}_1, \, \mathbf{i}_2, \, \dots, \, \mathbf{i}_M] \begin{bmatrix} \mathbf{G}_1 \\ \mathbf{G}_2 \\ \vdots \\ \mathbf{G}_M \end{bmatrix}$$

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• **Type-2** codes: Partitioning of parity check matrix **H** into subcodes  $C_{\ell}$  with parity check matrix **H**<sub> $\ell$ </sub>



#### LDPC Block Codes



#### **LDPC Block Codes**



Tanner graph (3,6) regular LDPC code:

 $\mathbf{H} = \begin{bmatrix} 1 & 0 & 1 & 1 & 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 1 \\ 1 & 0 & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 \end{bmatrix}$ 

Graph sparsely connected



#### Coupling construction via unwrapping:



Convolutional code structure

[Costello, Dolecek, Fuja, Kliewer, Mitchell, Smarandache, 2014]

Coupling construction via unwrapping:



Resulting Tanner graph:



Convolutional code structure

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Resulting Tanner graph:

#### **Spatially Coupled LDPC Codes: Performance**



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How can we build good nested codes with spatially coupled LDPC codes?

#### **Nested Spatially Coupled LDPC Codes**

Protograph representation of a type-1 nested spatially coupled LDPC code ensemble for M=2



#### **Results**



**Example:** Distributed fronthaul compression for cloud radio access networks (CRANs) in 5G



[Park, Simeone, Sahin, Shamai, 2014]

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- In the following: Lossy compression based on spatially coupled low-density generator matrix (LDGM) codes
  - Low encoding and decoding complexity (linear in time)
  - Performance very close to the rate-distortion limit

## Source Compression with Channel Codes

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 Idea: Treat source sequence as noisy codeword from some fictitious channel code (here a spatially coupled LDGM code)

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- Source encoding via modified belief propagation algorithm (channel decoding), windowed encoding for low latency
- Source decoding via channel encoding

#### **Coupling of Low-Density Generator Matrix Codes**



#### **Results: Symmetric Bernoulli Source**

![](_page_34_Figure_1.jpeg)

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![](_page_35_Figure_1.jpeg)

#### Take Aways

- Leveraging network gains in canonical multi terminal problems by nested SC-LDPC codes
  - relaying, broadcast, and cooperative diversity scenarios
- Low-complexity lossy and lossless compression with SC-LDGM codes

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- Leveraging network gains in canonical multi terminal problems by nested SC-LDPC codes
  - relaying, broadcast, and cooperative diversity scenarios
- Low-complexity lossy and lossless compression with SC-LDGM codes
- Example applications which can benefit from network compression gains:
  - Distributed compression for CRANs in 5G
  - Distributed compression of phasor measurement units in wide area measurement systems

#### • Open:

- ► Communication problem: Design of nested codes for *M*>2
- Compression problem: Design of nested codes and universal codes

## Follow Up...

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